# The universal property of $\mathcal{D}^-(\mathcal{A})$ based on Section 1.3.3 of Higher Algebra by Jacob Lurie

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# Conventions

Contrary to the conventions of Higher Algebra, we will omit the nerve construction of 1-categories from our notation.

We will often decorate functor categories with some symbols to mean a full subcategory spanned by functors satisfying some property. Examples:

- Fun<sup>rex</sup>: subcategory of right exact functors (i.e. those which preserve finite colimits).
- Fun<sup>|-|,II</sup>: subcategory of functors which preserve geometric realizations and finite coproducts.

We fix an abelian category A with enough projective objects. Changing universes if necessary, we assume that A is small.

## The main theorem

#### Definition (1.3.3.1)

Let C and C' be stable  $\infty$ -categories with t-structures. A functor  $F: C \to C'$  is called *right t-exact (rtex)* if it is exact (i.e. preserves finite limits and colimits) and  $F(C_{\geq 0}) \subseteq C'_{\geq 0}$ .

## Theorem (1.3.3.2)

Let C be a stable  $\infty$ -category equipped with a left complete t-structure.

Then

$$\begin{split} \mathsf{Fun}^{\textit{rtex},\mathcal{A}_{\textit{proj}}\mapsto\mathcal{C}^\heartsuit}(\mathcal{D}^-(\mathcal{A}),\mathcal{C}) \to \mathsf{Fun}^{\textit{rex}_{ab}}(\mathcal{A},\mathcal{C}^\heartsuit) \\ F \mapsto \tau_{\leqslant 0} \circ (F|\mathcal{D}^-(\mathcal{A})^\heartsuit) \end{split}$$

is an equivalence.

## A proof sketch for the main theorem

$$\begin{aligned} \mathsf{Fun}^{\mathsf{rtex},\mathcal{A}_{\mathsf{proj}}\mapsto\mathcal{C}^\heartsuit}(\mathcal{D}^-(\mathcal{A}),\mathcal{C}) \\ & \stackrel{\texttt{1.3.3.11}}{\textcircled{\hspace{0.5mm}}} \uparrow \\ \mathsf{Fun}^{|-|,\mathrm{II},\mathcal{A}_{\mathsf{proj}}\mapsto\mathcal{C}^\heartsuit}(\mathcal{D}^-_{\geqslant 0}(\mathcal{A}),\mathcal{C}_{\geqslant 0}) & \stackrel{\texttt{1.3.3.8}}{\longleftrightarrow} \mathsf{Fun}^{\mathrm{II}}(\mathcal{A}_{\mathsf{proj}},\mathcal{C}^\heartsuit) \\ & \stackrel{\uparrow^{1.3.3.9}}{\overset{\texttt{\scale}}{\downarrow}} \\ & \mathsf{Fun}^{\mathsf{rex}_{\mathsf{ab}}}(\mathcal{A},\mathcal{C}^\heartsuit) \end{aligned}$$

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## A lemma about geometric realizations

#### Lemma (1.3.3.10)

- 1. If an  $\infty$ -category C admits finite coproducts and geometric realizations of simplicial objects, then it admits all finite colimits. The converse is true if C is an n-category for some  $n \in \mathbb{N}$ .
- If F: C → D is a functor between ∞-categories admitting finite coproducts and geometric realizations of simplicial objects that preserves finite coproducts and geometric realizations, then F is right exact. The converse is true if C and D are n-categories for some n ∈ N.

# Proof of the geometric realization lemma

#### Proof sketch.

In order to construct all finite colimits from finite coproducts and geometric realizations, it is enough to construct coequalizers. Note there is a (non-full) inclusion  $\iota: ([1] \rightrightarrows [0]) \hookrightarrow \Delta^{\operatorname{op}}$  of the coequalizer shape into  $\Delta^{\operatorname{op}}$ . Moreover, using the pointwise formula for Kan extensions, one can show that  $\iota_1$  exists because finite coproducts exist.

Hence  $\operatorname{colim}_{[1]\rightrightarrows[0]} \simeq \operatorname{colim}_{\Delta^{\operatorname{op}}} \circ \iota_!$  exists.

For the part about *n*-categories, one shows that  $\operatorname{colim}_{\Delta^{\operatorname{op}}} \simeq \operatorname{colim}_{\Delta^{\operatorname{op}}_{\leqslant n}} \circ \iota_{\leqslant n}^*$  using some connectivity estimates after applying a Yoneda embedding and notes that the latter is a finite colimit.

# From right t-exact functors to right exact functors

## Lemma (1.3.3.11)

Let  $\mathcal{C},\,\mathcal{C}'$  be stable  $\infty\text{-categories}$  equipped with t-structures. Then

- 1. If C is right bounded, then restriction along  $C_{\geq 0} \hookrightarrow C$  induces an equivalence  $\operatorname{Fun}^{\operatorname{rtex}}(C, C') \simeq \operatorname{Fun}^{\operatorname{rex}}(C_{\geq 0}, C'_{\geq 0})$ .
- If C and C' are left complete, then C≥0 and C'≥0 admit geometric realizations of simplicial objects, and a functor F: C≥0 → C'≥0 is right exact if and only if it preserves finite coproducts and geometric realizations.

#### Proof the "rtex to rex" lemma, part 1

Proof sketch for (1).

Since  $C = \bigcup_n C_{\geq -n}$  by right boundedness, we have

 $\mathsf{Fun}^{\mathsf{rtex}}(\mathcal{C},\mathcal{C}')\simeq\mathsf{lim}(\ldots\to\mathsf{Fun}^{\mathsf{rex}}(\mathcal{C}_{\geqslant-1},\mathcal{C}'_{\geqslant-1})\to\mathsf{Fun}^{\mathsf{rex}}(\mathcal{C}_{\geqslant0},\mathcal{C}'_{\geqslant0})).$ 

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Now the tower on the RHS is essentially constant as  $\operatorname{Fun}^{rex}(\mathcal{C}_{\geq -n-1}, \mathcal{C}'_{\geq -n-1}) \to \operatorname{Fun}^{rex}(\mathcal{C}_{\geq n}, \mathcal{C}'_{\geq n})$  has a (homotopy) inverse given by "conjugation by  $\Sigma$ ".

# From right t-exact functors to right exact functors

## Lemma (1.3.3.11)

Let  $\mathcal{C},\,\mathcal{C}'$  be stable  $\infty\text{-categories}$  equipped with t-structures. Then

- 1. If C is right bounded, then restriction along  $C_{\geq 0} \hookrightarrow C$  induces an equivalence  $\operatorname{Fun}^{\operatorname{rtex}}(C, C') \simeq \operatorname{Fun}^{\operatorname{rex}}(C_{\geq 0}, C'_{\geq 0})$ .
- If C and C' are left complete, then C≥0 and C'≥0 admit geometric realizations of simplicial objects, and a functor F: C≥0 → C'≥0 is right exact if and only if it preserves finite coproducts and geometric realizations.

## Proof the "rtex to rex" lemma, part 2

#### Proof sketch for (2).

Left completeness means that

$$\mathcal{C}_{\geqslant 0} \simeq \mathsf{lim}(\ldots \to (\mathcal{C}_{\geqslant 0})_{\leqslant 1} \to (\mathcal{C}_{\geqslant 0})_{\leqslant 0}).$$

Moreover, each  $(C_{\geq 0})_{\leq n}$  is an finitely cocomplete (n + 1)-category. Hence, by  $\land$  Lemma 1.3.3.10 about geometric realizations, they all admit geometric realizations. Those are preserved by the truncation functors in the tower above, so C admits geometric realizations too.

The statement about preservation of colimits follows by a similar argument reducing the statement to the case of *k*-categories by virtue of the description of  $C_{\geq 0}$  as a limit of such.

# The t-structure on $\mathcal{D}^-(\mathcal{A})$

### Proposition (1.3.3.16)

The standard t-structure on  $\mathcal{D}^-(\mathcal{A})$  is right bounded and left complete.

#### Remark

Left completeness is ultimately reduced to the convergence of Postnikov towers of spaces by embedding  $\mathcal{D}^{-}(\mathcal{A})$  into the derived category of a presheaf category which we will be described in the next slide.

# A model for Ind-objects

## Proposition (1.3.3.13)

 $Ind(\mathcal{A})$  can be identified with  $\mathcal{A}^{\wedge} := Fun^{\times}(\mathcal{A}_{proj}^{op}, Set)$  which is again an abelian category with enough projective objects.

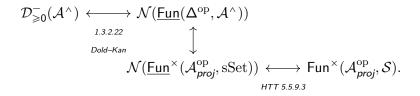
## Remark (1.3.3.15)

 $\mathcal{D}^{-}(\mathcal{A})$  can be identified with the full subcategory of  $\mathcal{D}^{-}(\mathcal{A}^{\wedge})$  consisting of objects whose homology belongs to (the Yoneda image of)  $\mathcal{A}$ .

## Product preserving presheaves

#### Proposition (1.3.3.14)

We have equivalences



Moreover, the restriction of the composite along the inclusion  $\mathcal{A}_{\text{proj}} \hookrightarrow \mathcal{D}^{-}(\mathcal{A}^{\wedge})$  corresponds to the Yoneda embedding.

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## A proof sketch for the main theorem

$$\begin{aligned} \mathsf{Fun}^{\mathsf{rtex},\mathcal{A}_{\mathsf{proj}}\mapsto\mathcal{C}^\heartsuit}(\mathcal{D}^-(\mathcal{A}),\mathcal{C}) \\ & \stackrel{\texttt{1.3.3.11}}{\textcircled{\hspace{0.5mm}}} \uparrow \\ \mathsf{Fun}^{|-|,\mathrm{II},\mathcal{A}_{\mathsf{proj}}\mapsto\mathcal{C}^\heartsuit}(\mathcal{D}^-_{\geqslant 0}(\mathcal{A}),\mathcal{C}_{\geqslant 0}) & \stackrel{\texttt{1.3.3.8}}{\longleftrightarrow} \mathsf{Fun}^{\mathrm{II}}(\mathcal{A}_{\mathsf{proj}},\mathcal{C}^\heartsuit) \\ & \stackrel{\uparrow^{1.3.3.9}}{\overset{\texttt{\scale}}{\downarrow}} \\ & \mathsf{Fun}^{\mathsf{rex}_{\mathsf{ab}}}(\mathcal{A},\mathcal{C}^\heartsuit) \end{aligned}$$

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# A more precise characterization

We still need to prove:

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Theorem (1.3.3.8)
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Let  ${\mathcal C}$  be an  $\infty\text{-category}$  admitting geometric realizations of simplicial objects.

Then

- 1. The restriction functor  $\operatorname{Fun}^{|-|}(\mathcal{D}^{-}_{\geq 0}(\mathcal{A}), \mathcal{C}) \to \operatorname{Fun}(\mathcal{A}_{\operatorname{proj}}, \mathcal{C})$  is an equivalence.
- A functor F ∈ Fun<sup>|-|</sup>(D<sup>-</sup><sub>≥0</sub>(A), C) preserves finite coproducts if and only if its restriction to A<sub>proj</sub> does.

## One last lemma

Part 1 is a corollary of the following:

## Lemma (1.3.3.17)

The essential image of  $\mathcal{D}_{\geq 0}^{-}(\mathcal{A})$  in  $\operatorname{Fun}^{\times}(\mathcal{A}_{proj}^{\operatorname{op}}, \mathcal{S})$  is the smallest full subcategory of  $\operatorname{Fun}(\mathcal{A}_{proj}^{\operatorname{op}}, \mathcal{S})$  that contains the image of the Yoneda embedding and is closed under geometric realizations.

#### Proof sketch.

The essential image contains the image of the Yoneda embedding. • Lemma 1.3.3.16 and • Lemma 1.3.3.11 imply that it is closed under

geometric realizations.

Moreover, (the image of) every object  $X \in \mathcal{D}_{\geq 0}^{-}(\mathcal{A})$  is equivalent to the geometric realization of the simplicial object in  $\mathcal{A}_{\text{proj}}$  that corresponds to (the chain complex underlying) X under the Dold–Kan correspondence.

# A more precise characterization

## Theorem (1.3.3.8)

Let  $\mathcal C$  be an  $\infty$ -category admitting geometric realizations of simplicial objects.

Then

- 1. The restriction functor  $\operatorname{Fun}^{|-|}(\mathcal{D}^{-}_{\geq 0}(\mathcal{A}), \mathcal{C}) \to \operatorname{Fun}(\mathcal{A}_{proj}, \mathcal{C})$  is an equivalence.
- A functor F ∈ Fun<sup>|-|</sup>(D<sup>-</sup><sub>≥0</sub>(A), C) preserves finite coproducts if and only if its restriction to A<sub>proj</sub> does.

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## Proof of part 2 of the more precise characterization

#### Proof sketch.

Let  $F: \mathcal{D}_{\geq 0}^{-}(\mathcal{A}) \to \mathcal{C}$  be a functor that preserves geometric realizations such that  $F' := F|\mathcal{A}_{\text{proj}}$  preserves finite coproducts. We need to show that F preserves finite coproducts. By possibly "extending"  $\mathcal{C}$  by virtue of HTT 5.3.5.7, we may assume that it has all colimits. Then HTT 5.5.8.15 says that F', as a

functor which preserves finite coproducts, extends to a functor in  $\mathsf{Fun}^{\mathsf{cocont}}(\mathsf{Fun}^{\times}(\mathcal{A}_{\mathsf{nroi}}^{\mathrm{op}}, \mathcal{S}), \mathcal{C}).$ 

Now restricting back to  $\mathcal{D}_{\geq 0}^{-}(\mathcal{A}) \subseteq \mathcal{D}_{\geq 0}^{-}(\mathcal{A}^{\wedge}) \simeq \operatorname{Fun}^{\times}(\mathcal{A}_{\operatorname{proj}}^{\operatorname{op}}, \mathcal{S})$ , we see that *F* itself also preserves finite coproducts.  $\Box$ 

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## A proof sketch for the main theorem

$$\begin{aligned} \mathsf{Fun}^{\mathsf{rtex},\mathcal{A}_{\mathsf{proj}}\mapsto\mathcal{C}^\heartsuit}(\mathcal{D}^-(\mathcal{A}),\mathcal{C}) \\ & \stackrel{\texttt{1.3.3.11}}{\textcircled{\hspace{0.5mm}}} \uparrow \\ \mathsf{Fun}^{|-|,\mathrm{II},\mathcal{A}_{\mathsf{proj}}\mapsto\mathcal{C}^\heartsuit}(\mathcal{D}^-_{\geqslant 0}(\mathcal{A}),\mathcal{C}_{\geqslant 0}) & \stackrel{\texttt{1.3.3.8}}{\longleftrightarrow} \mathsf{Fun}^{\mathrm{II}}(\mathcal{A}_{\mathsf{proj}},\mathcal{C}^\heartsuit) \\ & \stackrel{\uparrow^{1.3.3.9}}{\overset{\texttt{\scale}}{\downarrow}} \\ & \mathsf{Fun}^{\mathsf{rex}_{\mathsf{ab}}}(\mathcal{A},\mathcal{C}^\heartsuit) \end{aligned}$$

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